

Test 2 Answer

① a)

		q	1-q
		w	e
p	w	(2,2)	(2,3)
1-p	e	(3,2)	(0,0)

$$2q + 2(1-q) = 3q + 0(1-q)$$

$$2 = 3q$$

$$q = \frac{2}{3}$$

$$\text{Similarly, } p = \frac{2}{3}$$

Both watch with prob. $\frac{2}{3}$, eat with prob. $\frac{1}{3}$

b) Let $\sigma = pw + (1-p)e$ with $p = \frac{2}{3}$.

Let $\tau = qw + (1-q)e$

$\pi_{\sigma\sigma} = \pi_{\tau\sigma}$ for every τ .

$$\pi_{\sigma\sigma} - \pi_{\tau\sigma} = pq \cdot 2 + p(1-q) \cdot 2 + (1-p)q \cdot 3 + (1-p)(1-q) \cdot 0 \\ - [qq \cdot 2 + q(1-q) \cdot 2 + (1-q)q \cdot 3 + (1-q)(1-q) \cdot 0]$$

$$= (p-q)(2q + 2(1-q) - 3q)$$

$$= (p-q)(2-3q) = 3(p-q)\left(\frac{2}{3}-q\right) = 3\left(\frac{2}{3}-q\right)^2 > 0$$

for $q \neq \frac{2}{3}$.

Therefore σ is evolutionarily stable.

② a)

		q_1	q_2	q_3
		w	e	m
p_1	w	(1,1)	(0,2)	(0,2)
p_2	e	(2,0)	(-1,-1)	(2,0)
p_3	m	(2,0)	(0,2)	(-1,-1)

The following are equal:

$$\begin{cases} q_1 \cdot 1 + q_2 \cdot 0 + q_3 \cdot 0 \\ q_1 \cdot 2 + q_2 \cdot (-1) + q_3 \cdot 2 \\ q_1 \cdot 2 + q_2 \cdot 0 + q_3 \cdot (-1) \end{cases}$$

$$\Rightarrow \begin{cases} q_1 = 2q_1 - q_2 + 2q_3 \\ q_1 = 2q_1 - q_3 \end{cases} \Rightarrow \begin{cases} q_1 - q_2 + 2q_3 = 0 \\ q_1 - q_3 = 0 \end{cases}$$

Also: $q_1 + q_2 + q_3 = 1$

$$\Rightarrow q_1 = q_3 = \frac{1}{5}, q_2 = \frac{3}{5}$$

Both birds use w and m $\frac{1}{5}$ of the time, e $\frac{3}{5}$ of the time.

b)

		w	e	m
w	w	(1,1)	(0,2)	(0,2)
w	e	(2,0)	(-1,-1)	(2,0)
w	m	(2,0)	(0,2)	(-1,-1)

$$q \cdot -1 + (1-q) \cdot 2 = q \cdot 0 + (1-q) \cdot -1$$

$$-q + 2 - 2q = -1 + q$$

$$3 = 4q$$

$$q = \frac{3}{4}$$

Both birds use $\sigma = \frac{3}{4}e + \frac{1}{4}m$

$$\pi_1(e, \sigma) = \pi_1(m, \sigma) = -\frac{1}{4}$$

$$\pi_1(w, \sigma) = 0$$

Since $0 > -\frac{1}{4}$, (σ, σ) is not a Nash equilibrium.

- ③ σ^* = Worker's trigger strategy
 τ^* = Boss's trigger strategy

Let σ be another strategy for Worker.

(1) If Worker always works hard, payoffs are unchanged.

(2) If Worker first slacks off in round k ,

$$\left. \begin{array}{l} \text{Worker's payoff from round } k \text{ on} \\ \text{is at most} \end{array} \right\} \begin{array}{l} 5s^k + 3s^{k+1} + 3s^{k+2} + \dots \\ = 5s^k + 3s^k \cdot \frac{1}{1-s} \end{array}$$

$$\left. \begin{array}{l} \text{If Worker had continued to use } \sigma^* \\ \text{payoff from round } k \text{ on would be} \end{array} \right\} \begin{array}{l} 4s^k + 4s^{k+1} + \dots \\ = 4s^k \cdot \frac{1}{1-s} \end{array}$$

$$4\delta^k \cdot \frac{1}{1-\delta} \geq 5\delta^k + 3\delta^{k+1} \cdot \frac{1}{1-\delta} \Leftrightarrow$$

$$\frac{4}{1-\delta} \geq 5 + 3\delta \cdot \frac{1}{1-\delta} \Leftrightarrow 4 \geq 5(1-\delta) + 3\delta \Leftrightarrow$$

$$2\delta \geq 1 \Leftrightarrow \delta \geq \frac{1}{2}$$

Let τ be another strategy for Boss.

- (1) If Boss always pays H, payoffs are unchanged
- (2) If Boss first pays L in round k ,

Boss's payoff from round k on } $4\delta^k + 0\delta^{k+1} + 0\delta^{k+2} + \dots$
 is at most } $= 4\delta^k$

If Boss had continued to use τ^* , } $2\delta^k + 2\delta^{k+1} + \dots$
 payoff from round k on would be } $= 2\delta^k \cdot \frac{1}{1-\delta}$

$$2\delta^k \cdot \frac{1}{1-\delta} \geq 4\delta^k \Leftrightarrow \frac{2}{1-\delta} \geq 4 \Leftrightarrow \delta \geq \frac{1}{2}$$

So (σ^*, τ^*) is a Nash equilibrium for $\delta \geq \frac{1}{2}$.

- (4) a) Tough Guy always gets payoff of 2 from Steak and -2 from Quiche.
- b) Wimp's payoffs:

		HH	HL	LH	LL
Wimp	S	-8	-8	-2	-2
	Q	-4	2	-4	2

(*)

Bully encounters Tough Guy $\frac{1}{2}$ the time. His payoffs are

		HH	HL	LH	LL
Wimp	S	-2	-2	0	0
	Q	-2	-2	0	0

(Reason: In this case, Wimp's strategy is irrelevant, and Tough Guy always eats steak.)

Bully encounters Wimp $\frac{1}{2}$ the time. His payoffs are

		HH	HL	LH	LL
Wimp	S	2	2	0	0
	Q	2	0	2	0

Bully's payoffs are $\frac{1}{2}$ first matrix + $\frac{1}{2}$ second matrix:

		HH	HL	LH	LL
Wimp	S	0	0	0	0
	Q	0	-1	1	0

(**)

Answer from (*) and (**):

		HH	HL	LH	LL
Wimp	S	(-8,0)	(-8,0)	(-2,0)	(-2,0)
	Q	(-4,0)	(2,-1)	(-4,1)	(2,0)